### The intersection property on *k*-orbit polytopes

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# Maniplex

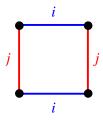
### Definition

Let  $\mathcal{M}$  be connected simple graph with a proper edge coloring  $c : \mathcal{M} \to \{0, \ldots, n-1\}$ . We say that  $\mathcal{M}$  is an *n*-maniplex if whenever |i - j| > 1, all paths of length 4 that alternate colors between i and j are closed.

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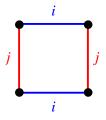


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Examples of maniplexes are the flag graphs of polytopes and maps, hence the vertices of a maniplex are often refered to as flags.

# Polytopality of maniplexes

### Definition

A maniplex satisfies the path intersection property (PIP) if whenever two vertices  $\Phi$  and  $\Psi$  are connected by some path with colors in a set I and another path with colors in J, then they are connected by a path with colors in I  $\cap$  J.

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#### Theorem

Garza-Vargas and Hubard, 2018 A maniplex  $\mathcal{M}$  is the flag graph of a polytope if and only if it satisfies the PIP.

#### Definition

Let  $\mathcal{M}$  be a maniplex and let  $\Gamma(\mathcal{M})$  be its automorphism group. Let  $H < \Gamma(\mathcal{M})$ . We define the symmetry type graph (STG) of  $\mathcal{M}$ with respect to H to be the quotient  $\mathcal{T}(\mathcal{M}, H) := \mathcal{M}/H$ .

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If there's an edge of color *i* between  $\Phi$  and  $\Psi$  on  $\Rightarrow$  there's an edge of color *i* between their orbits. If they are on the same orbit we draw a semi-edge on that orbit (not a loop).

The STG of a maniplex is a "non-simple maniplex", that is:

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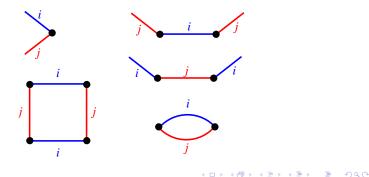
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Connected components of the graph induced by edges of colors colors i and j are one of the following:





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Given a graph X, its fundamental groupoid is the set of its reduced paths together with the (partial) operation "concatenation+reduction". It is denoted by  $\Pi(X)$ .

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#### Definition

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#### Definition

Let X be a graph. A voltage assignment is a (groupoid) anti-homomorphism  $\xi : \Pi(X) \to G$  where G is a group. In this case, G is called the voltage group.

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We only need to define the voltage of the arcs of X to define the voltage of all its paths. The voltage of a path is the product of the voltages of its arcs in reverse order.

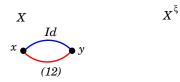
Given a graph X with a voltage assignment  $\xi$ , we may construct a "bigger" derived graph  $X^{\xi}$  in the following way:

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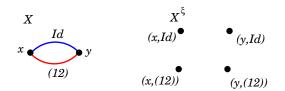
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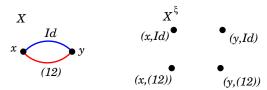
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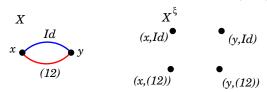
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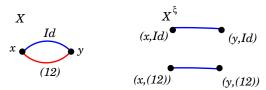
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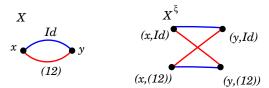
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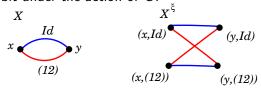
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We can recover X from  $X^{\xi}$  by identifying the vertices that are on the same orbit under the action of G.

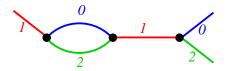
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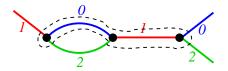
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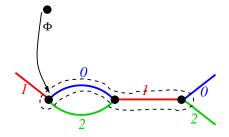


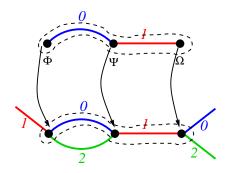
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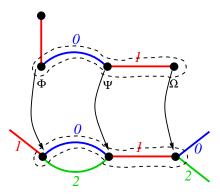


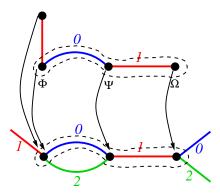
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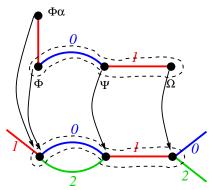




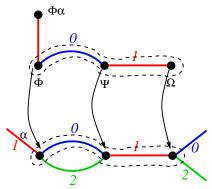


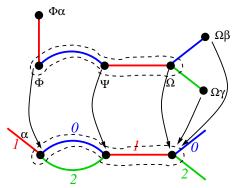


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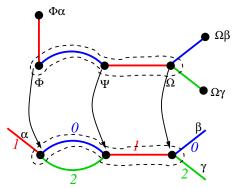


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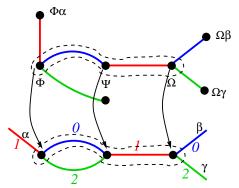


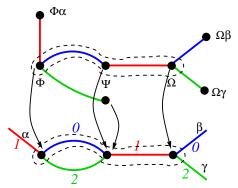


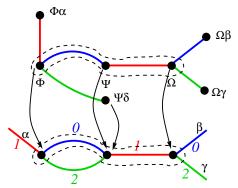
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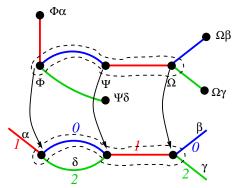
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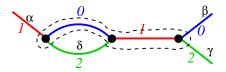




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- X<sup>ξ</sup> is simple iff all semi-edges have non-trivial voltage and all pairs of paralel edges have different voltages.
- X<sup>ξ</sup> is a maniplex iff both previous conditions hold, and whenever |i − j| > 1, the voltage of any path of length 4 that alternates colors between i and j is trivial.

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#### Definition

Let X be a non-simple maniplex. Let x and y be vertices and let I be a set of colors. We define  $\Pi_I^{x,y}(X)$  as the set of reduced paths from x to y that use colors in I.

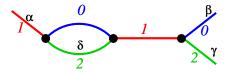
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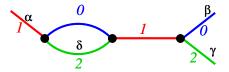
#### Theorem

Let X be a non-simple maniplex and  $\xi : \Pi(X) \to G$  a voltage assignment such that  $X^{\xi}$  is a maniplex. Then  $X^{\xi}$  is the flag graph of a polytope if and only if for every two vertices x and y and any two sets of colors I and J the following condition holds:

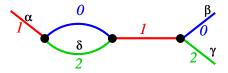
$$\xi(\Pi_I^{x,y}) \cap \xi(\Pi_J^{x,y}) = \xi(\Pi_{I \cap J}^{x,y}).$$





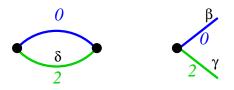


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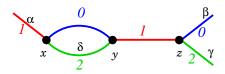
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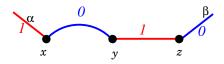
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• Maniplex:  $\delta^2 = 1$  and  $(\beta \gamma)^2 = 1$ .

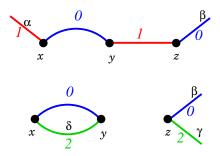




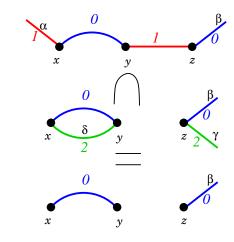




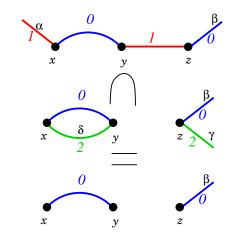
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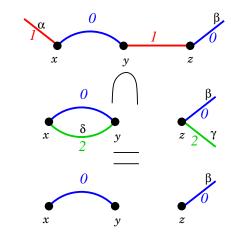


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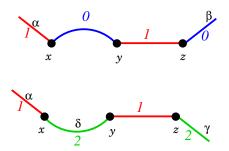
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- $\xi(\Pi_{\{0,1\}}^{\mathsf{x}}) = \langle \alpha, \beta \rangle.$   $\xi(\Pi_{\{0,2\}}^{\mathsf{x}}) = \langle \delta \rangle.$  $\xi(\Pi_{\{0\}}^{\mathsf{x}}) = 1.$
- $\blacktriangleright :: \langle \alpha, \beta \rangle \cap \langle \delta \rangle = 1.$



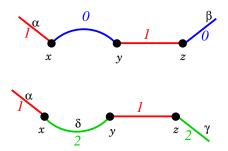
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 $\begin{aligned} & \xi(\Pi^{z}_{\{0,1\}}) = \langle \alpha, \beta \rangle. \\ & \xi(\Pi^{z}_{\{0,2\}}) = \langle \beta, \gamma \rangle. \\ & \xi(\Pi^{z}_{\{0\}}) = \langle \beta \rangle. \\ & \xi(\Pi^{z}_{\{0\}}) = \langle \beta, \gamma \rangle = \langle \beta \rangle. \end{aligned}$ 



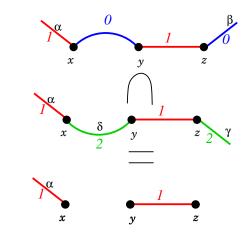
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• 
$$\xi(\Pi_{\{0,1\}}^x) = \langle \alpha, \beta \rangle.$$

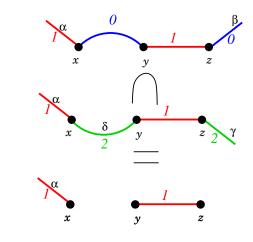


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• 
$$\xi(\Pi_{\{0,1\}}^{\mathsf{x}}) = \langle \alpha, \beta \rangle.$$
  
•  $\xi(\Pi_{\{1,2\}}^{\mathsf{x}}) = \langle \alpha, \gamma^{\delta} \rangle.$ 

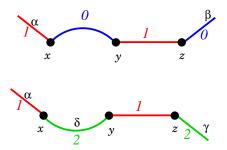


• 
$$\xi(\Pi_{\{0,1\}}^{\times}) = \langle \alpha, \beta \rangle.$$
  
•  $\xi(\Pi_{\{1,2\}}^{\times}) = \langle \alpha, \gamma^{\delta} \rangle.$   
•  $\xi(\Pi_{\{1\}}^{\times}) = \langle \alpha \rangle.$ 



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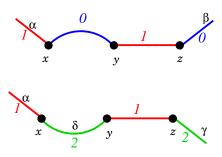
► 
$$\xi(\Pi_{\{0,1\}}^{x}) = \langle \alpha, \beta \rangle.$$
  
►  $\xi(\Pi_{\{1,2\}}^{x}) = \langle \alpha, \gamma^{\delta} \rangle.$   
►  $\xi(\Pi_{\{1\}}^{x}) = \langle \alpha \rangle.$   
►  $\therefore \langle \alpha, \beta \rangle \cap \langle \alpha, \gamma^{\delta} \rangle = \langle \alpha \rangle.$ 



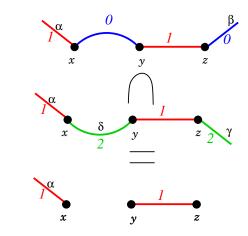
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• 
$$\xi(\Pi_{\{0,1\}}^{x,y}) = \langle \alpha, \beta \rangle.$$

## • $\xi(\Pi_{\{0,1\}}^{x,y}) = \langle \alpha, \beta \rangle.$ • $\xi(\Pi_{\{1,2\}}^{x,y}) = \delta \langle \alpha, \gamma^{\delta} \rangle.$

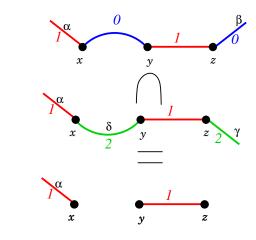


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• 
$$\xi(\Pi_{\{0,1\}}^{x,y}) = \langle \alpha, \beta \rangle.$$

$$\xi(\Pi_{\{0,1\}}) = \langle \alpha, \beta \rangle.$$
  
 
$$\xi(\Pi_{\{1,2\}}^{x,y}) = \delta \langle \alpha, \gamma^{\delta} \rangle.$$



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- $\xi(\Pi_{\{0,1\}}^{x,y}) = \langle \alpha, \beta \rangle.$
- $\xi(\prod_{\{1,2\}}^{x,y}) = \delta\langle \alpha, \gamma^{\delta} \rangle.$
- $\xi(\Pi_{\{1\}}^{\hat{\mathbf{x}},\hat{\mathbf{y}}}) = \emptyset.$  $\therefore \langle \alpha, \beta \rangle \cap \delta \langle \alpha, \gamma^{\delta} \rangle = \emptyset.$

Using this method we get these seven conditions (and many other trivial ones):

- $\blacktriangleright \ \langle \alpha, \beta \rangle \cap \langle \delta \rangle = 1$
- $\blacktriangleright \ \langle \alpha, \beta \rangle \cap \langle \beta, \gamma \rangle = \langle \beta \rangle$
- $\blacktriangleright \langle \alpha, \beta \rangle \cap \langle \alpha, \gamma^{\delta} \rangle = \langle \alpha \rangle$
- $\blacktriangleright \langle \alpha, \beta \rangle \cap \langle \alpha^{\delta}, \gamma \rangle = 1$

$$\blacktriangleright \langle \alpha, \beta \rangle \cap \delta \langle \alpha, \gamma^{\delta} \rangle = \emptyset$$

$$\blacktriangleright \langle \delta \rangle \cap \langle \alpha, \gamma^{\delta} \rangle = 1$$

$$\blacktriangleright \ \langle \beta, \gamma \rangle \cap \langle \alpha^{\delta}, \gamma \rangle = \langle \gamma \rangle$$

Using this method we get these seven conditions (and many other trivial ones):

- $\blacktriangleright \ \langle \alpha, \beta \rangle \cap \langle \delta \rangle = 1$
- $\blacktriangleright \ \langle \alpha, \beta \rangle \cap \langle \beta, \gamma \rangle = \langle \beta \rangle$
- $\blacktriangleright \langle \alpha, \beta \rangle \cap \langle \alpha, \gamma^{\delta} \rangle = \langle \alpha \rangle$
- $\blacktriangleright \ \langle \alpha,\beta\rangle \cap \langle \alpha^{\delta},\gamma\rangle = 1$

$$\blacktriangleright \langle \alpha, \beta \rangle \cap \delta \langle \alpha, \gamma^{\delta} \rangle = \emptyset$$

$$\blacktriangleright \langle \delta \rangle \cap \langle \alpha, \gamma^{\delta} \rangle = 1$$

$$\blacktriangleright \ \langle \beta, \gamma \rangle \cap \langle \alpha^{\delta}, \gamma \rangle = \langle \gamma \rangle$$

#### Theorem

Let H be a group. There exists a polytope  $\mathcal{P}$  such sthat  $H < \Gamma(\mathcal{P})$ and  $\mathcal{T}(\mathcal{P}, H)$  is the graph from the example iff H is generated by four involutions  $\alpha, \beta, \gamma, \delta$  such that  $(\beta\gamma)^2 = 1$  and the previous intersection properties hold.

# Thank you!